Relational Query Languages: Taxonomy

Relational Query Languages		
Relational	Relational Calculus	
Algebra	Tuple Calculus	Domain Calculus
ISBL, SQL	QUEL, SQL	QBE,
		MS-Access

When a query language is **relationally complete**

Equivalence of the relational algebra, the tuple calculus, and the domain calculus

Relational Queries and Solutions - Relational Algebra - I

students(**sname**, gpa, fname) faculty(**fname**, office)

1. all advisors

 π_{fname} students

2. all student records with gpa's over 3.0

 $\sigma_{gpa > 3.0}$ students

3. all advisors of students with gpa's over 3.0

 $\pi_{\text{fname}} \sigma_{\text{gpa} > 3.0}$ students

4. offices of all advisors of students with gpa's over 3.0

 π_{office} (faculty * $\sigma_{\text{gpa} > 3.0}$ students)

Relational Queries and Solutions - Relational Algebra - II

students(**sname**. gpa, fname) offerings(**cno**, **sem**, fname) took(**sname**, **cno**, **sem**, grade)

1. students who took a course with their advisor

 π_{sname} (students * offerings * took)

2. students who never took a course with their advisor

 π_{sname} students - π_{sname} (students * offerings * took)

3. students who took a course their advisor never taught

 π_{sname} ($\pi_{\text{sname, cno}}$ took - $\pi_{\text{sname, cno}}$ (students * offerings))

3.5) students who took a course with someone other than their advisor

 $\pi_{\text{sname}}(\pi_{\text{sname,cno,sem}} \text{ took - } \pi_{\text{sname,cno,sem}}(\text{students*offerings}))$

4. students who only took courses with their advisor

 π_{sname} students - π_{sname} ($\pi_{\text{sname, cno, sem}}$ took -

 $\pi_{\text{sname, cno, sem}}$ (students * offerings))